

Modeling of Time in Discrete-Event Simulation of Systems-on-Chip

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Work partially supported by
HELP ANR project

MEMOCODE, July 2011

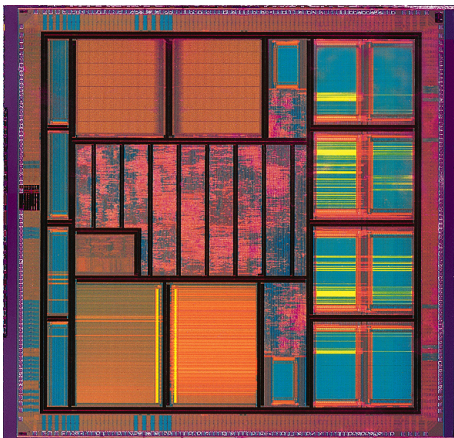
Outline

- 1 Transaction Level Modeling and jTLM
- 2 Time and Duration in jTLM
- 3 Applications
- 4 Implementation
- 5 Conclusion

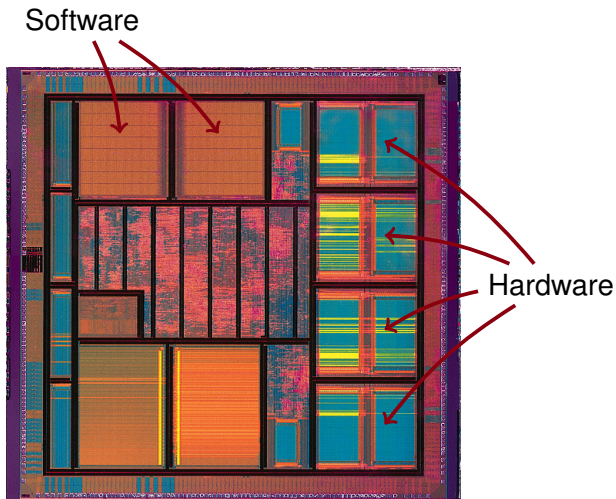
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Modern Systems-on-a-Chip



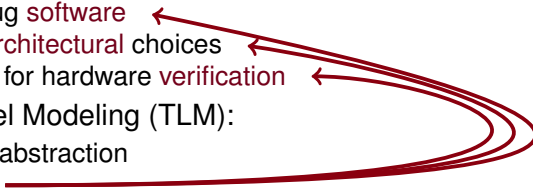
Modern Systems-on-a-Chip



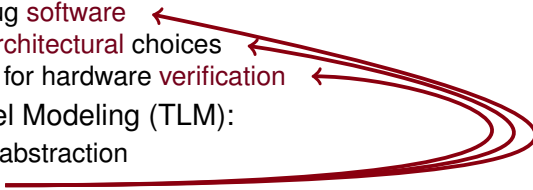
Transaction-Level Modeling

- (Fast) simulation essential in the design-flow
 - ▶ To write/debug **software**
 - ▶ To validate **architectural** choices
 - ▶ As reference for hardware **verification**

Transaction-Level Modeling

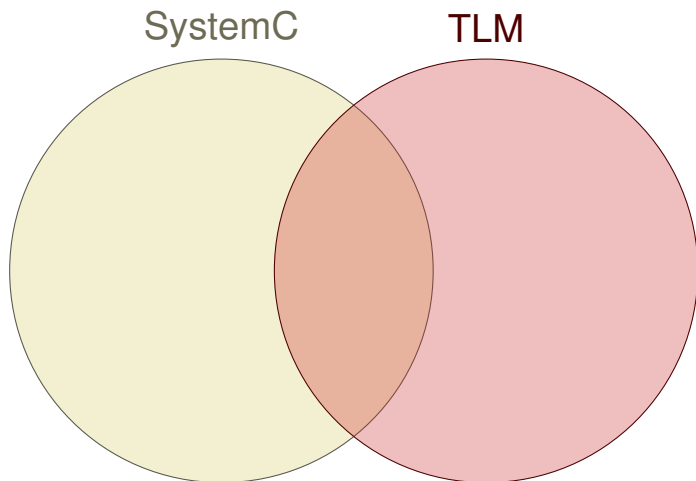
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 - Transaction-Level Modeling (TLM):
 - ▶ High level of abstraction
 - ▶ Suitable for
- 

Transaction-Level Modeling

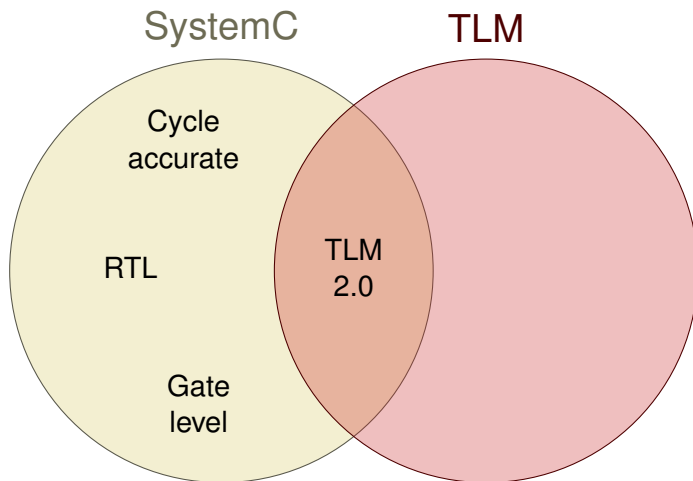
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Industry Standard = SystemC/TLM

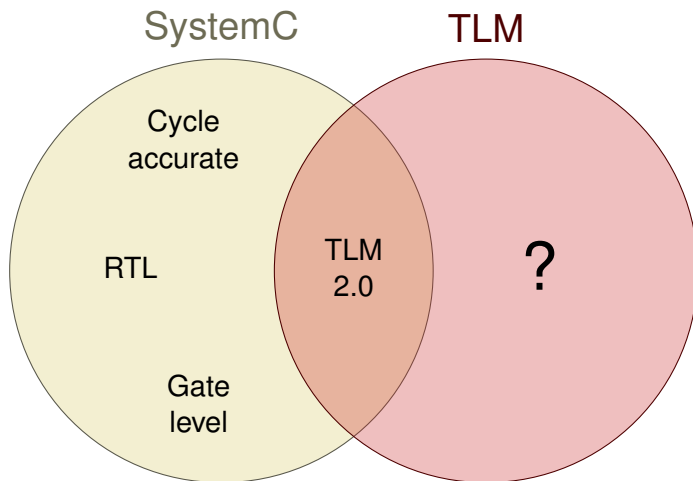
SystemC/TLM vs. "TLM Abstraction Level"



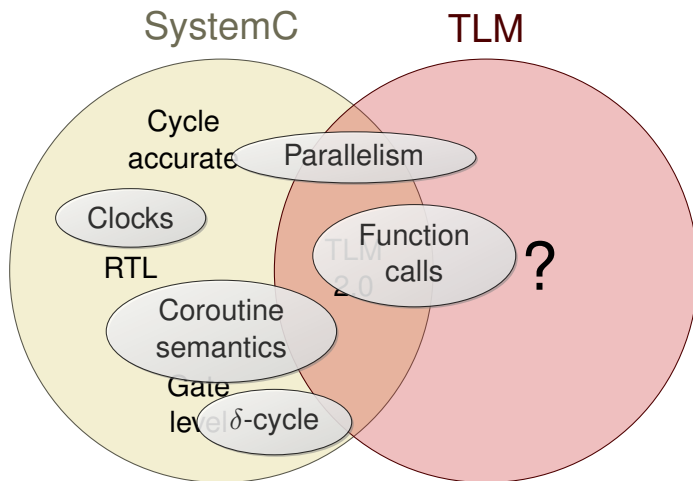
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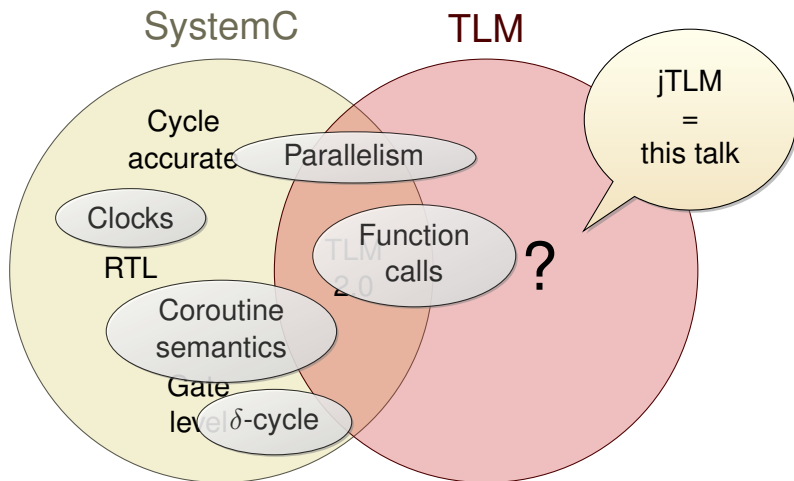
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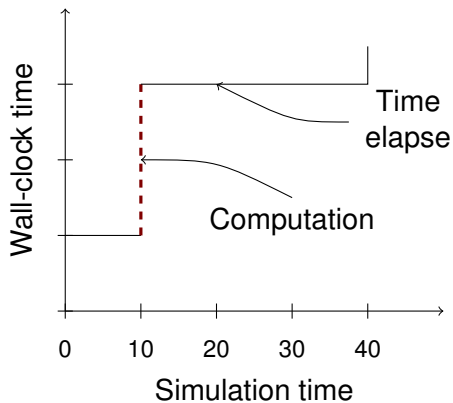
jTLM: goals and peculiarities

- jTLM's goal: define "TLM" independently of SystemC
 - ▶ **Not** cooperative (true parallelism)
 - ▶ **Not** C++ (Java)
 - ▶ **No** δ -cycle
- Interesting features
 - ▶ Small and simple code (\approx 500 LOC)
 - ▶ Nice experimentation platform
- Not meant for production

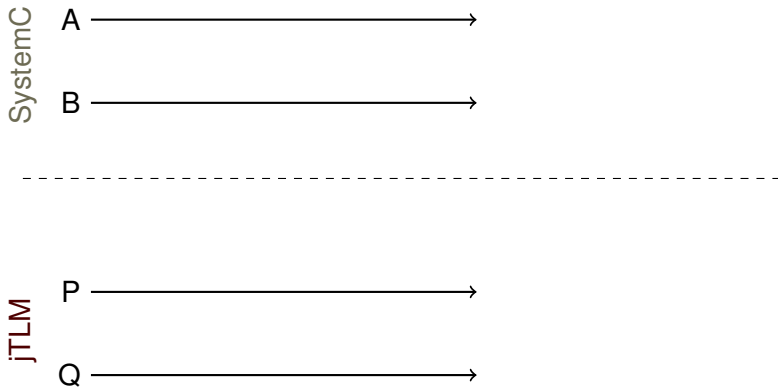
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Simulation Time Vs Wall-Clock Time

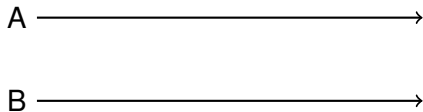


Time in SystemC and jTLM



Time in SystemC and jTLM

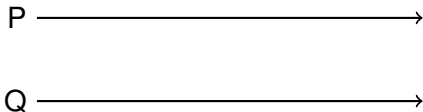
SystemC



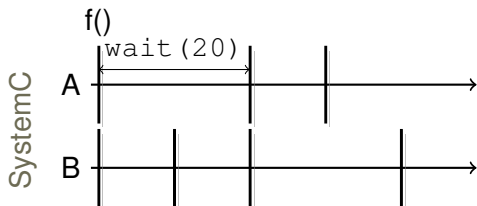
Process A:

```
//computation  
f();  
//time taken by f  
wait(20, SC_NS);
```

jTLM



Time in SystemC and jTLM



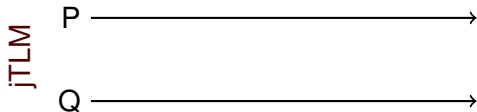
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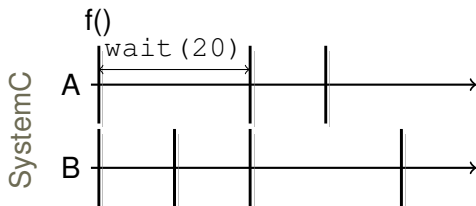
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f();
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```
//time taken by f
```

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wait(20, SC_NS);
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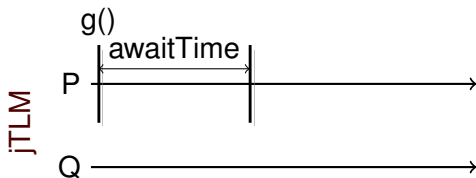


Time in SystemC and jTLM



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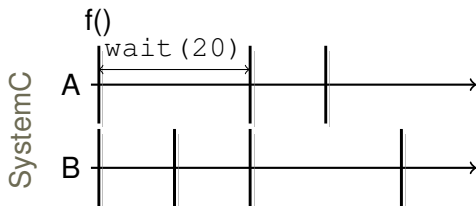
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Process P:

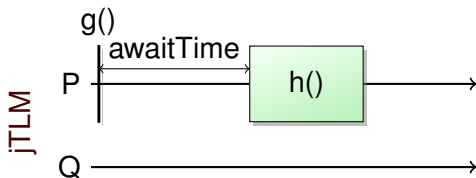
```
g();
awaitTime(20);
```

Time in SystemC and jTLM



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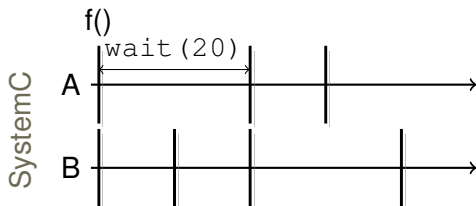
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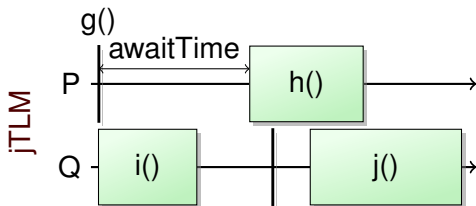
```
g();
awaitTime(20);
consumeTime(15) {
    h();
}
```

Time in SystemC and jTLM



Process A:

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f();
//time taken by f
wait(20, SC_NS);
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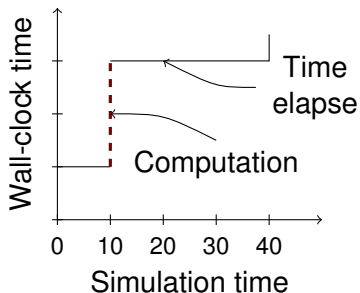


Process P:

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g();
awaitTime(20);
consumeTime(15) {
    h();
}
```

Time à la SystemC: `awaitTime(T)`

- By default, time does not pass
⇒ instantaneous tasks
- `awaitTime(T)` :
let other processes execute
for T time units



```
f(); // instantaneous
awaitTime(20);
```

Task with Known Duration: `consumesTime (T)`

- Semantics:

- ▶ Start and end dates known
- ▶ Actions contained in task spread in between

- Advantages:

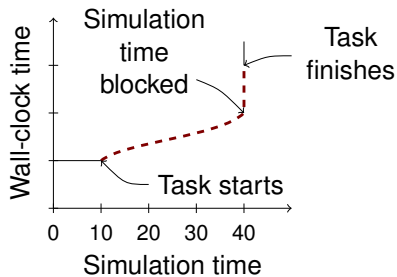
- ▶ Model closer to actual system
- ▶ Less bugs hidden
- ▶ Better parallelization



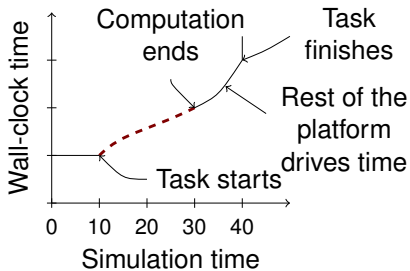
```
consumesTime (15) {  
    f1 ();  
    f2 ();  
    f3 ();  
}  
  
consumesTime (10) {  
    g ();  
}
```


Execution of `consumesTime (T)`

Slow computation



Fast computation



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Exposing Bugs

Example bug: mis-placed synchronization:

```
flag = true;           while(!flag)
awaitTime(5);          awaitTime(1);
writeIMG();           awaitTime(10);
awaitTime(10);        readIMG();
```

⇒ bug never seen in simulation

Exposing Bugs

Example bug: mis-placed synchronization:

```

flag = true;           while(!flag)
awaitTime(5);          awaitTime(1);
writeIMG();           ||   awaitTime(10);
awaitTime(10);        readIMG();

```

⇒ bug never seen in simulation

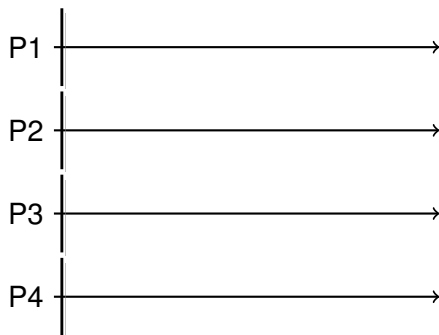
```

consumesTime(15) {    while(!flag)
    flag = true;      ||   awaitTime(1);
    writeIMG();       awaitTime(10);
}                     readIMG();

```

⇒ strictly more behaviors, including the buggy one

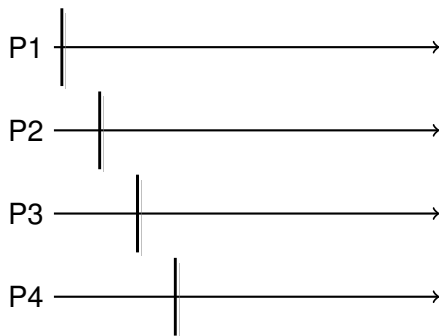
Parallelization



jTLM's Semantics

- Simultaneous tasks run
in parallel

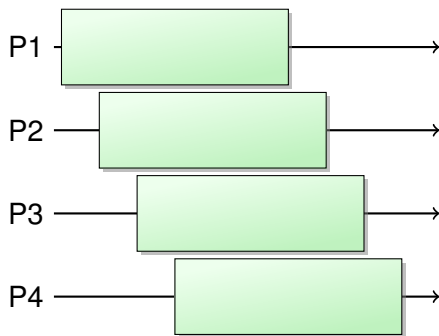
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jTLM's Semantics

- Simultaneous tasks run **in parallel**
- Non-simultaneous tasks don't

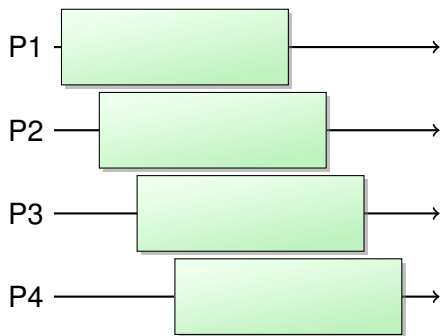
Parallelization



jTLM's Semantics

- Simultaneous tasks run **in parallel**
- Non-simultaneous tasks don't
- Overlapping tasks do

Parallelization



jTLM's Semantics

- Simultaneous tasks run **in parallel**
- Non-simultaneous tasks don't
- Overlapping tasks do

- **Back to SystemC:**

- ▶ Parallelizing within δ -cycle = great if you have clocks
- ▶ Simulation time is the bottleneck with quantitative/fuzzy time

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Time Queue and `awaitTime (T)`



Process P:

```
▶ f ();  
  awaitTime (50);
```

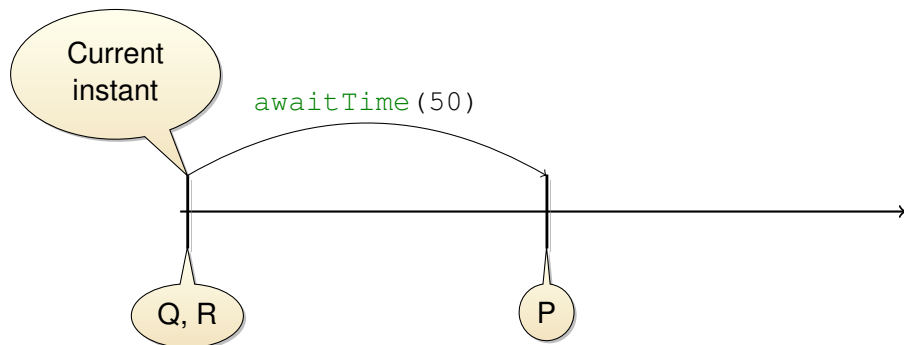
Process Q:

```
▶ h ();  
  awaitTime (30);  
  g ();  
  awaitTime (30);
```

Process R:

```
▶ i ();  
  awaitTime (90);
```

Time Queue and `awaitTime (T)`



Process P:

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f();
▷awaitTime (50);
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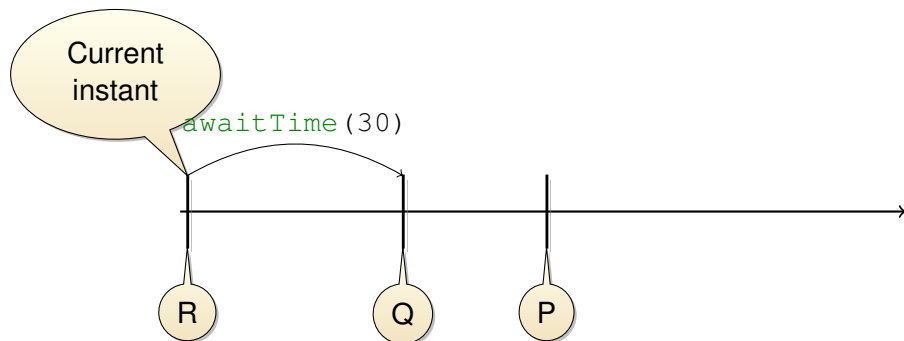
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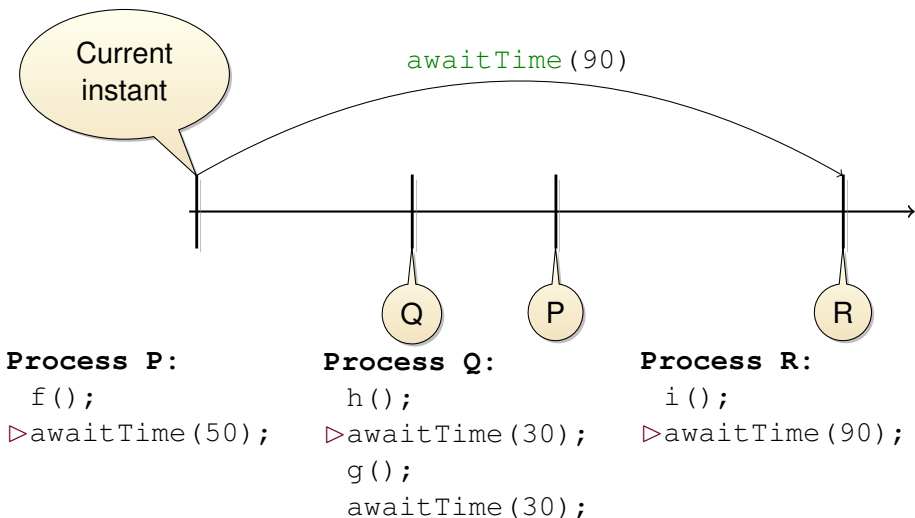
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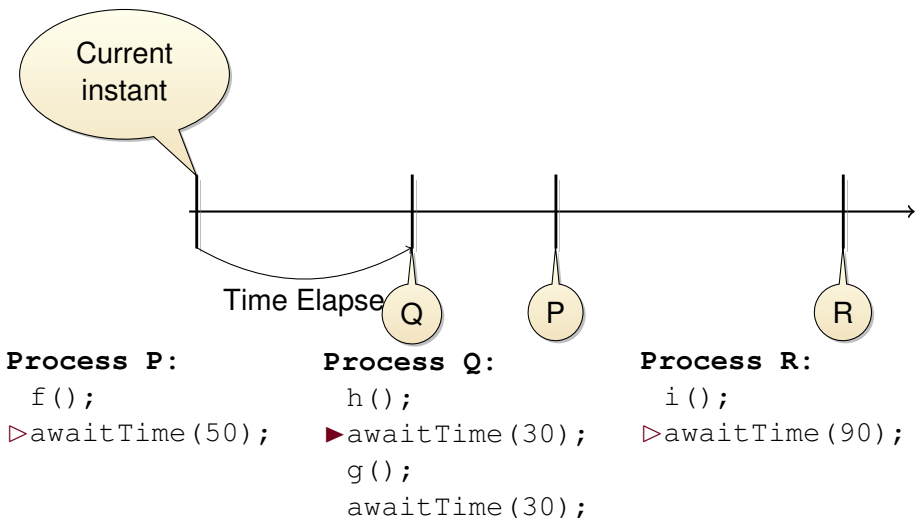
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```

```
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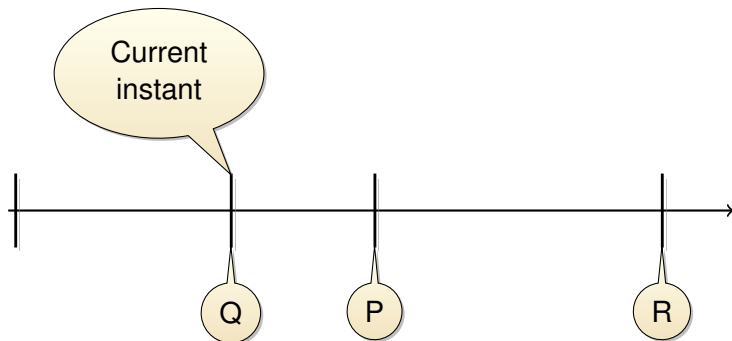
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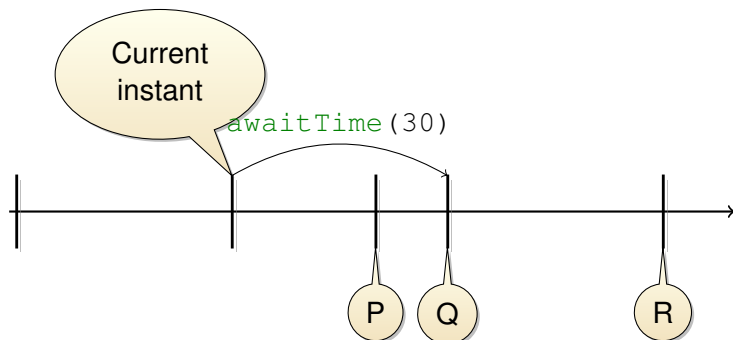
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Time Queue and `consumesTime (T)`

What about `consumesTime (T)` ?

Time Queue and `consumeTime (T)`

Current
instant

P, Q, R

Process P:

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  consumeTime (50) {
    g ();
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h ();
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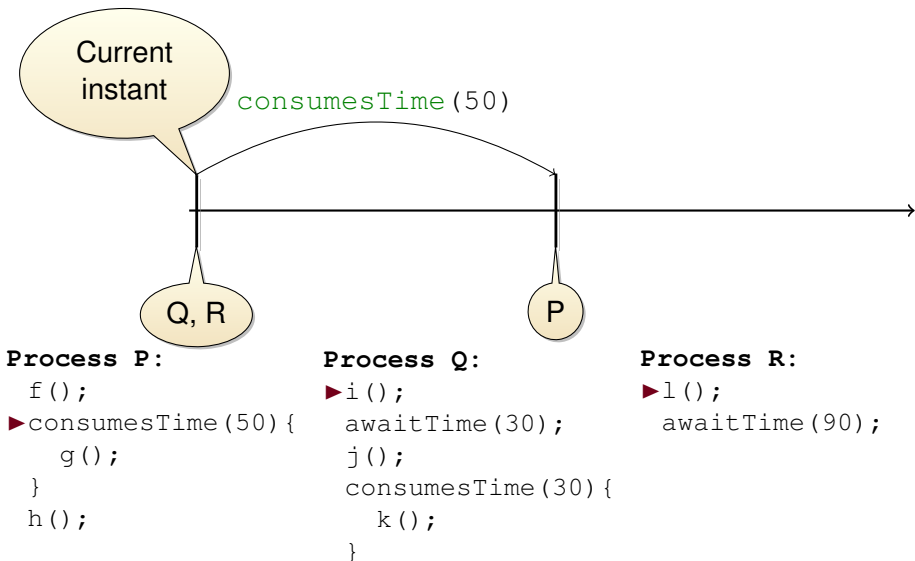
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  consumeTime (30) {
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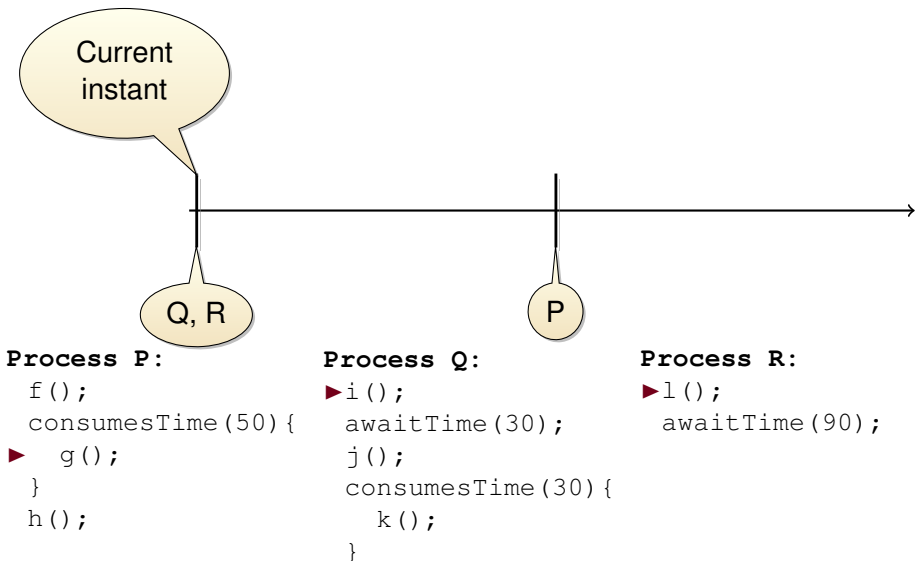
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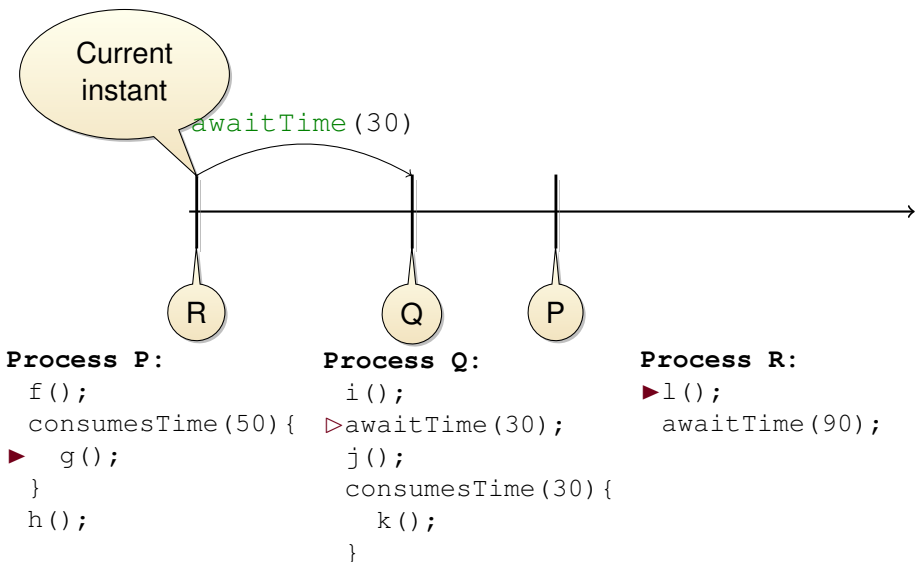
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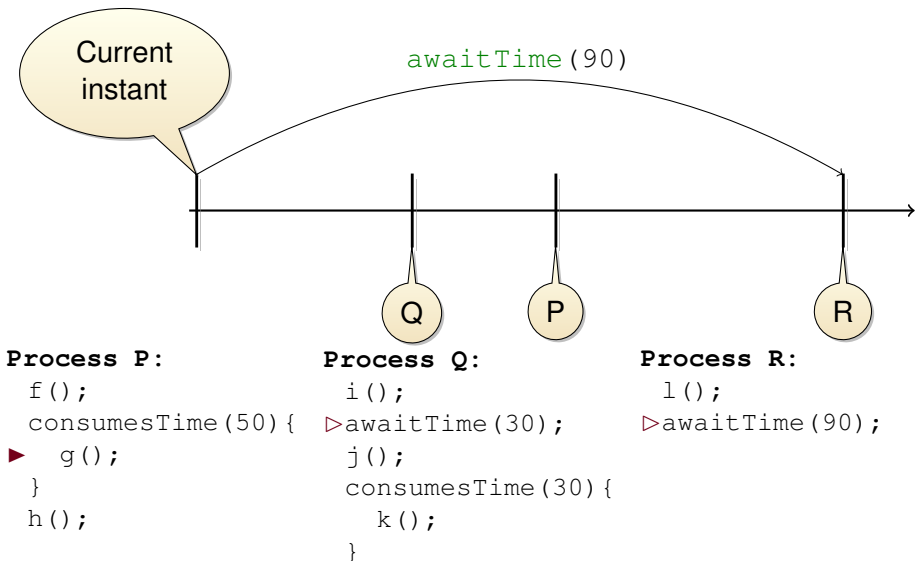
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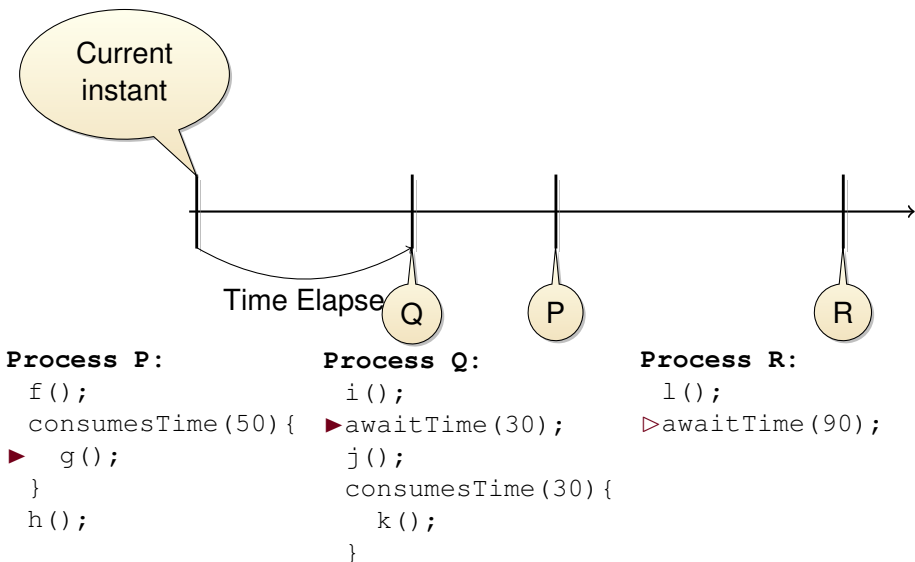
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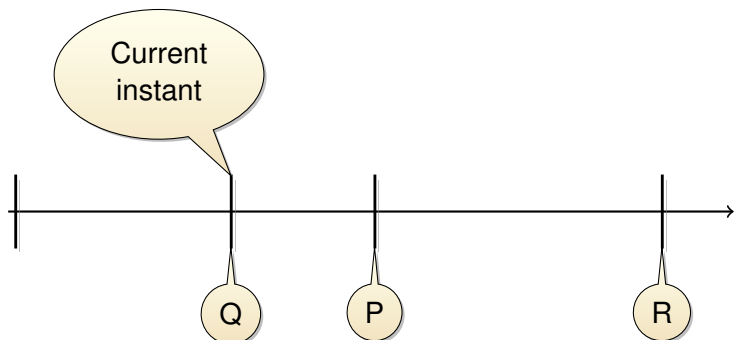
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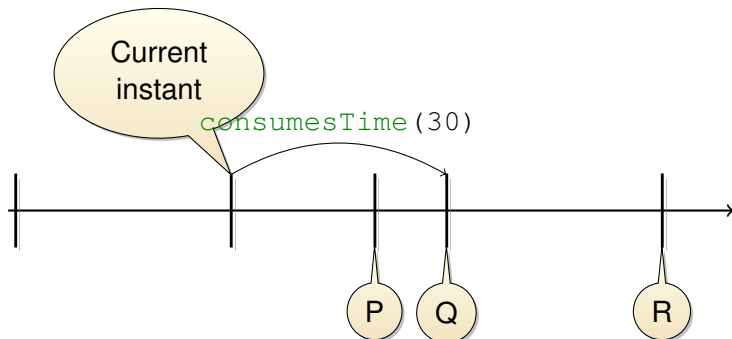
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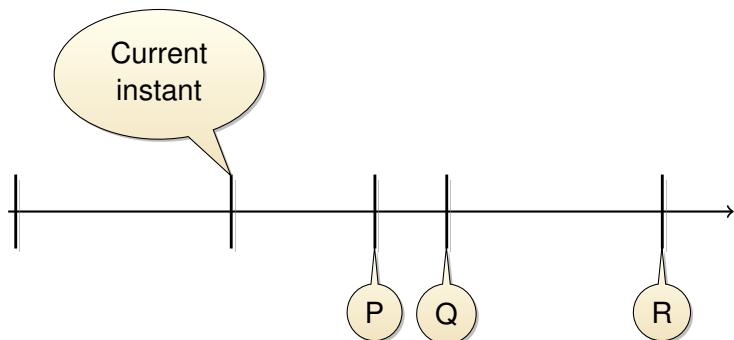
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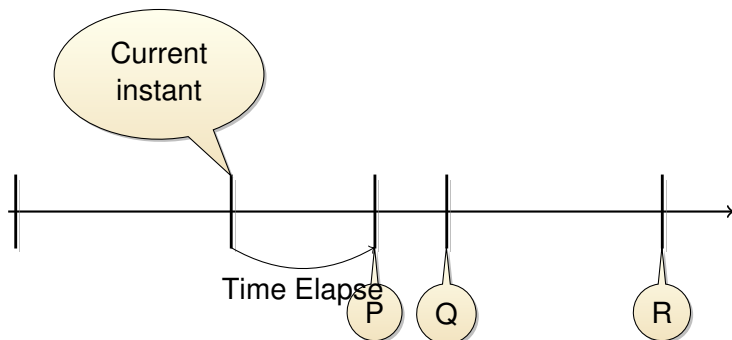
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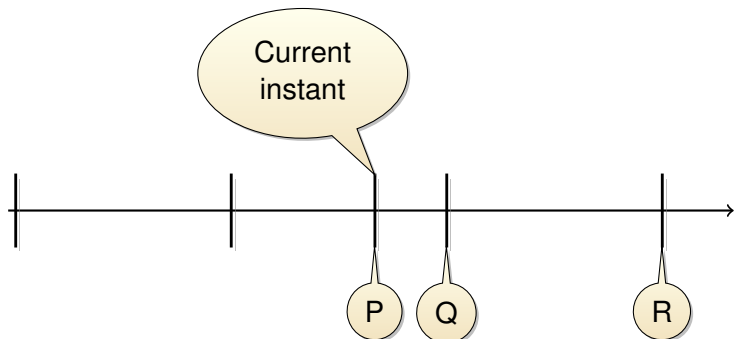
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Perspectives

- Summary
 - ▶ Tasks with duration
 - ▶ Exhibit more behaviors/bugs
 - ▶ Better parallelization
- Skipped from the talk (cf. paper)
 - ▶ Tasks with a priori unknown duration
 - ▶ jTLM's cooperative mode
- Perspectives
 - ▶ Adapt the ideas to SystemC (ongoing, not so hard)
 - ▶ Run-time Verification to explore schedules (science-fiction)
 - ▶ Open-Source Release?

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Thank you! ~> Questions?